# Dynamic Programming in Futhark 

Martin Elsman, DIKU, University of Copenhagen

October 23, 2021

This note describes how the Futhark [1, 2] dynamic programming library dpsolver ${ }^{1}$ can be used to find fixpoints for functions from $R^{n}$ to $R^{n}$ and how solutions to multiple instances of a dynamic programming problem can be computed in parallel on GPUs. We give both single-dimensional and multi-dimensional examples and we show how the Futhark automatic differentiation feature may relieve programmers from specifying explicitly the Jacobian matrices, which are necessary for using dpsolver's fast converging Newtonian functionality.

## Introduction

A standard approach for finding fixpoints for numerical functions from $R^{n}$ to $R^{n}$ is to use the technique of successive approximations. Following Section 4 of Numerical Dynamic Programming in Economics, by John Rust [3], the dynamic programming solver that we shall apply here first uses a number of successive approximation steps before it applies a more efficient Newtonian method for narrowing in on a fixpoint. The latter method requires that the user specifies how to compute the Jacobian matrix (of type $R^{n \times n}$ ) given an approximate fixpoint. The Jacobian is then computed for each Newtonian step.

## Example: Intersection of a circle and a quadratic equation

Following the example in Jim Lander's MAT 461/561 lecture notes, we first set out to find the intersection between the unit circle $\left(x_{1}^{2}+x_{2}^{2}=1\right)$ and the quadratic equation $x_{2}=x_{1}^{2}$. ${ }^{2}$
We first define an operator for which we want to find a fixpoint. To ensure that the natural matrix norm of the Jacobian matrix for the function is less than $1\left(0 \leq x_{1} \leq 1\right.$ and $\left.0 \leq x_{2} \leq 1\right)$, we give the following definition of the fixpoint operator $G$ :

$$
\begin{equation*}
G\left(x_{1}, x_{2}\right)=\left(\sqrt{x_{2}}, \sqrt{1-x_{1}^{2}}\right) \tag{1}
\end{equation*}
$$

Without having to define the Jacobian matrix for the function, we can find an approximation to the fixpoint using the successive approximation functionality of the dpsolver library.

[^0]We first import the library dpsolver and instantiate the contained module dpsolver to the f64 representation of floats:

```
import "dpsolver"
module dps = mk_dpsolver f64
```

The function dps.sa that we shall apply has the following type:

```
val sa [m] : (f:[m]t->[m]t) -> (v:[m]t) -> (p:param) -> (b:t)
    -> ([m]t, bool, i64, t, t)
```

Here $t$ is identical to $f 64$ due to the $f 64$ module instantiation of the mk_dpsolver parameterised module.

We now define the bellman equation for which we want to find a fixpoint:

```
let bellman (x1:f64) (x2:f64) : (f64,f64) =
    (f64.sqrt x2, f64.sqrt(1 - x1**2))
```

Before we can make use of the library module, we define a small utility function wrap2, which takes a function of type $f 64 \rightarrow>f 64 \rightarrow$ (f64, f64) and turns it into a function of type [2] f64 -> [2]f64, which is compatible with the successive approximation functionality in the dps module:

```
let wrap2 (f: f64 -> f64 -> (f64,f64)) (a:[2]f64) : [2]f64 =
    let (x,y) = f a[0] a[1]
    in [x,y]
```

The following Futhark entry point makes a call to the dps.sa function with the above bellman function wrapped as a parameter:

```
entry test_sa (sa_max:i64) (sa_tol:f64)
    : ([2]f64, bool, i64, f64) =
    let v0 = [0.5, 0.5]
    let ap = dps.default with sa_max = sa_max
                                    with sa_tol = sa_tol
    let (res, b, i, tol, _) = dps.sa (wrap2 bellman) v0 ap 0
    in (res, b, i, tol)
```

The function dps.sa also takes an initial approximation as argument (v0) together with an ap value that defines some slightly modified default parameter settings (max iterations, max tolerance, etc.)
We can now call the function:

```
> test_sa 60i64 1.0e-3f64
([0.7855137639650378f64, 0.6177294754734614f64], true, 56i64,
    8.769119278559945e-4f64)
```

We see that after 56 iterations, a fixpoint is found with a tolerance below $1 \mathrm{e}-3$, meaning that the last iteration step contributed to a change in value of less than $1 \mathrm{e}-3$ for both $x_{1}$ and $x_{2}$. For improved precision, many more iterations are required:

```
> test_sa 200i64 1.0e-9f64
([0.7861513784203592f64, 0.6180339890667096f64], true, 186i64,
    9.120002530949023e-10f64)
```


## Faster convergence with Newton's method

The function $G$, as defined in (1), has the following Jacobian matrix:

$$
J_{G}\left(x_{1}, x_{2}\right)=\left[\begin{array}{cc}
0 & \frac{1}{2 \sqrt{x_{2}}} \\
\frac{-x_{1}}{\sqrt{1-x_{1}^{2}}} & 0
\end{array}\right]
$$

The following version of the bellman function takes its input as an array of size 2 and returns, along with the function result, the Jacobian matrix, relative to the argument:

```
let bellman_j (a:[2]f64) : ([2]f64, [2][2]f64) =
    let x1 = a[0]
    let x2 = a[1]
    let res = [f64.sqrt x2, f64.sqrt(1-x1**2)]
    let j = [[0 , 1/(2*f64.sqrt x2) ],
            [-x1/(f64.sqrt(1-x1**2)) , 0 ]]
    in (res, j)
```

The function dps.poly that we shall apply has the following type:

```
val poly [m]: (f: [m]t -> ([m]t,[m][m]t)) -> (v:[m]t) -> (p:param)
    -> (b:t) -> ([m]t,[m][m]t,bool,i64,i64,i64,t)
```

Again, here $t$ is identical to $f 64$ due to the $f 64$ module instantiation of the mk_dpsolver parameterised module. The function finds a fixpoint for the function $f$ using a combination of successive approximation iterations and Newton-Kantorovich iterations. The initial guess is $v$ and the parameter $p$ controls the iteration passes. The function $f$ should return a pair of a new next approximation and the Jacobian matrix for the function $f$ relative to the argument given. The function returns a 7-tuple containing an approximate fixpoint, a Jacobian matrix for the fixpoint, a boolean specifying whether the algorithm converged (according to the values in p ), the number of iterations used for the total sa iterations, the total Newton-Kantoovich iterations, and the number of roundtrips. The 7 'th element of the result tuple is the tolerance of the last two fixpoint approximations (maximum of each dimension).

```
entry test_poly (sa_max:i64) : ([2]f64, bool, i64, i64, i64, f64) =
    let v0 = [0.5, 0.5]
    let ap = dps.default with sa_max = sa_max
    let (res, _, b, i, j, k, tol) = dps.poly bellman_j v0 ap 0
    in (res, b, i, j, k, tol)
> test_poly 5i64
([0.7861513777574233f64, 0.6180339887498949f64], true, 5i64, 4i64, 1i64,
    1.1102230246251565e-16f64)
```

Notice that the progammer has manually provided code for computing the Jacobian matrix for the function. The result is a fixpoint with a tolerance below $1 \mathrm{e}-15$, computed with an initial number of 5 successive approximation iterations followed by 4 Newtonial iterations ( 1 roundtrip was used).

## Futhark AD

We can relieve the programmer from manually providing the code for the Jacobian matrix by using the automatic differentiation feature of Futhark, which provides a function jvp that performs forward automatic differentiation on arbitrary Futhark functions (featured in the Futhark clean-ad branch, but not yet featured in the master branch). An alternative is to encode float computations using so-called dual-numbers, following the approach of AD with dual numbers, but we shall not dive into this possibility here.

We first define a function wrapj that takes a function of type [n] [m]. [n]f64-> [m]f64 and turns it into a function of type $[\mathrm{n}][\mathrm{m}] .[\mathrm{n}] \mathrm{f} 64->([\mathrm{m}] \mathrm{f} 64,[\mathrm{~m}][\mathrm{n}] \mathrm{f} 64)$ that, besides from the function result, returns the Jacobian matrix of the function:

```
let idd n i = tabulate n (\j -> if i==j then 1f64 else Of64)
let wrapj [n] [m] (f: [n]f64-> [m]f64) (x:[n]f64) : ([m]f64,[m][n]f64) =
    (f x, tabulate n (jvp f x <-< idd n) |> transpose)
```

Functions wrapped with the wrapj function can now be used directly with the dps.poly function. Let's try it out in practice:

```
entry test_poly_jvp (sa_max:i64) : ([2]f64,bool,i64,i64,i64,f64) =
    let v0 = [0.5, 0.5]
    let ap = dps.default with sa_max = sa_max
    let (res, _, b, i, j, k, tol) =
        dps.poly ((wrapj <-< wrap2) bellman) v0 ap 0
    in (res, b, i, j, k, tol)
> test_poly_jvp 5i64
([0.7861513777574233f64, 0.6180339887498949f64], true, 5i64, 4i64, 1i64,
    1.1102230246251565e-16f64)
```

We see that we get the same results with test_poly_jvp as we get with test_poly.

## Going parallel

The iterative approaches that the dpsolver functionality implements for finding fixpoints are inherently sequential, except from the matrix operations applied in the Newton-Kantorovich iterations (assuming a high-number of dimensions). Instead of parallelising the actual fixpoint resolution, we shall see how we can find many fixpoints in parallel, which is sometimes a useful approach for speeding up an application.

Following up on the task of finding intersection points between a circle and a simple quadratic equation, let us investigate how the x-dimension of the intersection points changes when the
circle radius increases.
We first parameterise the bellman equation over the radius of the circle:

```
let bellmanr (r:f64) (a:[2]f64) : ([2]f64, [2][2]f64) =
    let f (a:[2]f64) = [f64.sqrt a[1], f64.sqrt(r**2-a[0]**2)]
    let res = f a
    let j = [[0 , 1/(2*f64.sqrt a[1])],
                [-a[0]/f64.sqrt(r**2-a[0]**2) , 0 ]]
    in (res, j)
```

We then create an entry point that implements an outer map over a call to dps.poly with varying radius:

```
let linspace (n: i64) (start: f64) (end: f64) : [n]f64 =
    tabulate n (\i -> start + f64.i64 i * ((end-start)/f64.i64 n))
entry test_polyr (n:i64) (sa_max:i64)
    : (bool, i64, [n]f64, [n]f64) =
    let ap = dps.default with sa_max = sa_max
    let rs = linspace n 1 20
    let ress = map (\r -> let v0 = [0.5,0.5]
                let (res, _, b, i, j, _k, _tol) =
                        dps.poly (bellmanr r) v0 ap 0
                in (r,res[0],b,i+j)) rs
    let converged = reduce (&&) true (map (.2) ress)
    let xs = map (.1) ress
    let iterations = reduce (+) 0 (map (.3) ress)
    in (converged, iterations, rs, xs)
```

Here is a call to test_polyr with 4 different radius values (between 1 and 20) and an sa_max value of 3 :

```
> test_polyr 4i64 3i64
(true, 25i64, [1.0f64, 5.75f64, 10.5f64, 15.25f64],
    [0.7861513777574233f64, 2.2960178985163853f64, 3.164158343195599f64, 3.841639561393954f64])
```

We can use the plot functionality of literate Futhark to plot 1000 points relating radius values with associated found $x$-values (and compare it with a plot of the sqrt-function):

```
entry test_polyr_rxs (n:i64) (sa_max:i64)
    : ([n]f64, [n]f64) =
    test_polyr n sa_max |> (\(_,_,rs,xs) -> (rs,xs))
let xys f n start end =
    unzip (map (\x -> (x, f x)) (linspace n start end))
entry sqrt_coords = xys f64.sqrt
> :plot2d {rxs=test_polyr_rxs 1000i64 3i64,
```



## A few single-dimensional examples

We now consider a single-dimensional case, for which we want to find the $x$ for which $f(x)=\cos x$.

```
entry test_poly1d (sa_max : i64)
    : ([1]f64, [1][1]f64, bool, i64, i64, i64, f64) =
    let ap = dps.default with sa_max = sa_max
    in dps.poly (\x -> ([f64.cos x[0]],
                                    [[- f64.sin x[0]]]))
            [0.7] ap 0
> test_poly1d 0i64
([0.7390851332151607f64], [[-0.6736120230211678f64]], true, 0i64, 3i64, 1i64,
0.0f64)
```

For another example, we want to compute $\sqrt{2}$ by finding the fixpoint to the equation $f(x)=$ $\frac{1}{2}\left(x+\frac{2}{x}\right)$.
entry test_sqrt (sa_max : i64)
: ([1]f64, [1][1]f64, bool, i64, i64, i64, f64) =
let ap = dps.default with sa_max = sa_max
in dps.poly ( $\backslash \mathrm{x}->([0.5 *(x[0]+2 / x[0])]$,
[[ 2*x[0] ]] )
) [1.4] ap 0
> test_sqrt 0i64
([1.414213562373095f64], [[2.828427124746191f64]], true, 0i64, 4i64, 1i64,
1.1102230246251565e-15f64)

Remarkably, in 4 steps we reach a fixpoint of $1.41421356237 \ldots$ with a tolerance of $1.11 \mathrm{e}-15$.

## Conclusion

We have seen how we can use the dpsolver library to solve multi-dimensional fixpoint equations. We have also seen how we can solve multiple problems in parallel using Futhark's second-order array combinators.

## References

[1] Elsman, M., Henriksen, T. and Oancea, C.E. 2018. Parallel programming in Futhark. Department of Computer Science, University of Copenhagen.
[2] Henriksen, T., Serup, N.G.W., Elsman, M., Henglein, F. and Oancea, C.E. 2017. Futhark: Purely functional GPU-programming with nested parallelism and in-place array updates. Proceedings of the 38th ACM SIGPLAN conference on programming language design and implementation (New York, NY, USA, 2017), 556-571.
[3] Rust, J. 1996. Numerical dynamic programming in economics. Elsevier. 619-729.


[^0]:    ${ }^{1}$ The Futhark library dpsolver is based on a Matlab library implemented by Bertel Schjerning, ECON, University of Copenhagen.
    ${ }^{2}$ Analytically, the solution can easily be found by solving the quadratic equation $x_{2}^{2}+x_{2}-1=0$, which leads to the solution $x_{1}=0.786151377757$ and $x_{2}=0.61803398874989$.

